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DIAGRAM REWRITING AND OPERADS

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DIAGRAM REWRITING AND OPERADS

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Abstract. — We give a survey of a diagrammatic syntax for *PROs* and *PROPs*, which are related to the theory of operads and bialgebras. Using *diagram rewriting*, we obtain *presentations of PROs by generators and relations*. In some cases, we even get *convergent rewrite systems*.

Résumé (Réécriture de diagrammes et opérades). — Nous donnons un aperçu de la syntaxe diagrammatique pour les *PROs* et les *PROPs*, qui sont liés à la théorie des opérades et des bigèbres. En utilisant la *réécriture de diagrammes*, on obtient des *présentations de PROs par générateurs et relations*. Dans certains cas, on obtient même des *systèmes de réécriture convergents*.

Except for Sections 4 and 7, most of the material presented in this paper comes from [12], which was inspired by [2].

1. PROs and PROPs

Definition 1. — A PRO (or product category) is a strict monoidal category, that is a (small) category $\mathbf{C}$ equipped with some associative functor $*$: $\mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$ and a unit object, such that the set of objects of $\mathbf{C}$ is $\mathbb{N}$, and $p * q = p + q$ for all $p, q \in \mathbb{N}$. In particular, the unit object is 0.

In order to define a PRO, since objects are already known, it suffices to give the set $\mathbf{C}(p, q)$ of morphisms $f : p \rightarrow q$ for all $p, q \in \mathbb{N}$, together with:

- a sequential composition $g \circ f : p \rightarrow r$ for any $f : p \rightarrow q$ and $g : q \rightarrow r$;
- a parallel composition $f * f' : p + p' \rightarrow q + q'$ for any $f : p \rightarrow q$ and $f' : p' \rightarrow q'$;
- an identity $\text{id}_p : p \rightarrow p$ for all $p \in \mathbb{N}$.

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This terminology will be clear in the next section. Of course, those two compositions must be associative, with units:

- $(h \circ g) \circ f = h \circ (g \circ f)$ for any $f : p \rightarrow q$, $g : q \rightarrow r$, and $h : r \rightarrow s$;
- $(f * f') * f'' = f * (f' * f'')$ for any $f : p \rightarrow q$, $f' : p' \rightarrow q'$, and $f'' : p'' \rightarrow q''$;
- $f \circ \text{id}_p = f = \text{id}_q \circ f$ and $f * \text{id}_0 = f = \text{id}_0 * f$ for any $f : p \rightarrow q$.

But they must also be compatible (law of *interchange*):

- $(g \circ f) * (g' \circ f') = (g * g') \circ (f * f')$ for any $f : p \rightarrow q$, $g : q \rightarrow r$, $f' : p' \rightarrow q'$, and $g' : q' \rightarrow r'$;
- $\text{id}_p * \text{id}_q = \text{id}_{p+q}$ for all $p, q \in \mathbb{N}$.

Here are typical examples:

- the PRO $\mathfrak{F}$, where a morphism $f : p \rightarrow q$ is a map from $\{1, \dots, p\}$ to $\{1, \dots, q\}$;
- the PRO $\mathfrak{M} \subset \mathfrak{F}$, where a morphism $f : p \rightarrow q$ is a monotone map from $\{1, \dots, p\}$ to $\{1, \dots, q\}$;
- the PRO $\mathbf{L}(\mathbb{K})$, where a morphism $f : p \rightarrow q$ is a $\mathbb{K}$ -linear map from $\mathbb{K}^p$ to $\mathbb{K}^q$ (or a $q \times p$ matrix) for any commutative field $\mathbb{K}$.

Compositions are obvious:

- $\circ$ is composition of maps (or product of matrices);
- $*$ is disjoint union (for $\mathfrak{F}$), ordered sum (for $\mathfrak{M}$), or direct sum (for $\mathbf{L}(\mathbb{K})$).

If we remove the object 0 from the PRO $\mathfrak{M}$, then we get the *simplicial category* Δ .

Definition 2. — A PRO $\mathbf{C}$ is reversible if all $\mathbf{C}(p, p)$ are groups and $\mathbf{C}(p, q) = \emptyset$ whenever $p \neq q$.

In order to define such a PRO, it suffices to give a group $\mathbf{C}_p = \mathbf{C}(p, p)$ for all p , together with a parallel composition $f * g \in \mathbf{C}_{p+q}$ defined for any $f \in \mathbf{C}_p$ and $g \in \mathbf{C}_q$. Note that a reversible PRO is a groupoid, but the condition $\mathbf{C}(p, q) = \emptyset$ for $p \neq q$ is not necessary to get a groupoid. Here are typical examples:

- the PRO $\mathfrak{S} \subset \mathfrak{F}$, where $\mathfrak{S}_p$ is the p -th *symmetric group*;
- the PRO $\mathfrak{B}$, where $\mathfrak{B}_p$ is the p -th *braid group*;
- the PRO $\mathbf{GL}(\mathbb{K}) \subset \mathbf{L}(\mathbb{K})$, where $\mathbf{GL}_p(\mathbb{K})$ is the p -th *linear group over $\mathbb{K}$* ;
- the PRO $\mathbf{O} \subset \mathbf{GL}(\mathbb{R})$, where $\mathbf{O}_p \subset \mathbf{GL}_p(\mathbb{R})$ is the p -th *orthogonal group*.

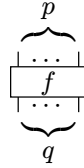
Definition 3. — A PROP (or product and permutation category) is a PRO $\mathbf{C} \supset \mathfrak{S}$.

For instance, both $\mathfrak{F}$ and $\mathbf{L}(\mathbb{K})$ are PROPs, but not $\mathfrak{M}$. PROPs are introduced in [16], with a slightly different definition, but of course, our notion is equivalent.

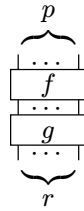
2. Diagrams

We recall the *diagrammatic syntax* of [12]:

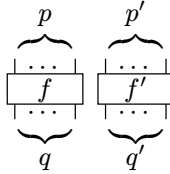
- a morphism $f : p \rightarrow q$ is pictured as a box with p inputs and q outputs:



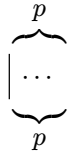
- for $f : p \rightarrow q$ and $g : q \rightarrow r$, the sequential composition $g \circ f : p \rightarrow r$ is pictured as follows:



- for $f : p \rightarrow q$ and $f' : p' \rightarrow q'$, the parallel composition $f * f' : p + p' \rightarrow q + q'$ is pictured as follows:



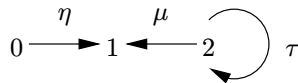
- the identity $\text{id}_p : p \rightarrow p$ is pictured as follows:



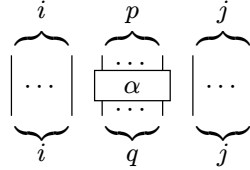
- in particular, $\text{id}_0 : 0 \rightarrow 0$ is pictured as an empty diagram.

Definition 4. — A signature is a graph $\mathcal{S}$ with vertices in $\mathbb{N}$. An edge $\alpha : p \rightarrow q$ in $\mathcal{S}$ is called a symbol with p inputs and q outputs.

For instance, the following signature will be introduced in the next section:



Definition 5. — An elementary diagram built over signature $\mathcal{S}$ is a formal parallel composition $\text{id}_i * \alpha * \text{id}_j : i + p + j \rightarrow i + q + j$, where $\alpha : p \rightarrow q$ is a symbol of $\mathcal{S}$, and $i, j \in \mathbb{N}$. It is pictured as follows:

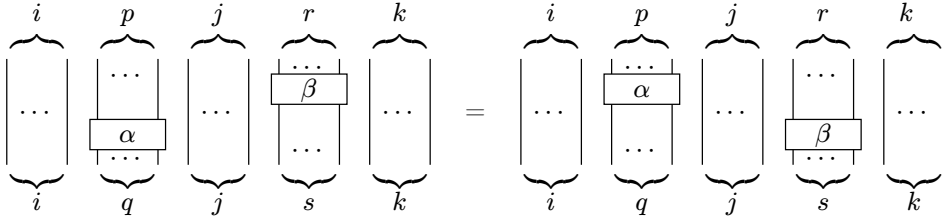


Definition 6. — A diagram built over signature $\mathcal{S}$ is a formal sequential composition $\phi_n \circ \dots \circ \phi_1 : p_0 \rightarrow p_n$, where $\phi_1 : p_0 \rightarrow p_1, \phi_2 : p_1 \rightarrow p_2, \dots, \phi_n : p_{n-1} \rightarrow p_n$ are elementary diagrams. In particular, we get $\text{id}_{p_0} : p_0 \rightarrow p_0$ when $n = 0$.

Definition 7. — The free PRO $\mathcal{S}^*$ consists of all diagrams built over a signature $\mathcal{S}$, modulo the commutation laws:

$$(\text{id}_i * \alpha * \text{id}_{j+s+k}) \circ (\text{id}_{i+p+j} * \beta * \text{id}_k) = (\text{id}_{i+q+j} * \beta * \text{id}_k) \circ (\text{id}_i * \alpha * \text{id}_{j+r+k})$$

for any symbols $\alpha : p \rightarrow q$ and $\beta : r \rightarrow s$, and for all $i, j, k \in \mathbb{N}$.



The commutation laws are necessary to get a PRO, which must satisfy interchange. In fact, a morphism of $\mathcal{S}^*$ can also be considered as a formal (sequential and parallel) composition of symbols modulo associativity, units, and interchange.

3. Presentations by generators and relations

Definition 8. — A relation $\rho = \sigma$ (over a signature $\mathcal{S}$) is given by two diagrams $\rho, \sigma : p \rightarrow q$ built over $\mathcal{S}$.

Definition 9. — A presentation of a PRO $\mathbf{C}$ consists of a signature $\mathcal{S}$ together with a set $\mathcal{R}$ of relations over $\mathcal{S}$, such that $\mathbf{C} \simeq \mathcal{S}^* / \leftrightarrow_{\mathcal{R}}^*$, where $\leftrightarrow_{\mathcal{R}}^*$ is the congruence generated by $\mathcal{R}$.